

Processor Analysis and Tuning

- **What is CPU utilization**
- **Theory of processor speed**
- **Common problems**
- **LPAR, HiperDispatch, Horizontal**
- **Overview of Processors**
- **Processor measurements**
- **Steal time**
- **Master Processor**
- **PLDV, Dispatch rates**
- **MFC, SMT**

Processor Performance Concepts - Utilization

What is important?

- TOTAL IFL Utilization
- LPAR Utilization
- “My” share

CPU Utilization used for:

- Performance Analysis
- Capacity Planning
- Accounting/Chargeback

Utilization measured in many ways

- Virtual Linux measures what? Percent of something.
- z/VM under LPAR measures what? CPU Seconds
- Hardware measurement only valid method of measuring CPU

Processor Performance Concepts - Utilization

What is “CPU Utilization”?

Percent of Percent misleading

- Can not be used directly for capacity planning
- Can not be used directly for accounting/chargeback
- Often misleading for performance analysis

All zVPS numbers are measured in CPU Seconds

- Percent is always based on CPU seconds divided by wall clock

Impacts measurements of

- LPAR
- z/VM Virtual Machines
- Linux processes
- zVSE Jobs/Partitions

Processor CAPACITY Concepts – Adding IFLs

Adding processors in MP environment reduces MSU / CP in z/OS

Adding processors in z/VM MP environment

- Reduces queuing time
- increases CPU available
- Cost money

Service levels comparison at 50% CPU Queue time

- One CPU: 50% busy, queue time = cpu time
- Two CPUs: 70% busy, queue time = cpu time
- Three CPUS: 79% busy, queue time = cpu time
- Four CPUs: 84% busy, queue time = cpu time

The 2nd CPU almost triples the capacity at target service level

The 3rd CPU adds 70% to capacity

More processors operate at higher utilization with good performance

- **Many expensive mistakes based on mis-understanding**
- **More hardware costs money, financial choices**
- **Understand Impacts (common mistakes)**
 - Number of vcpu in LPAR (impact on relative weight?)
 - Number of vcpu in virtual machine (impact on relative share)?
 - ALL single threaded large CPU Consumers! MP doesn't help?
 - ALL CPU intensive applications
- **CPU response time is a function of processor speed and CPUs.**
 - Faster processors mean shorter service time.
 - More processors means lowering queuing time.
- **Which provides better response time,**
 - 2 x 20 MIP engines, or 1 x 40 MIP engine?
 - at what percent utilization?

Processor Performance Concepts

- Which provides better response time multi-process workload,
 - 2 x 20 MIP engines, (**LESS QUEUEING**)
 - 1 40 MIP engine? (**SMALLER SERVICE TIME**)
- at what percent utilization? (“x” is cpu requirement)
 - MM1: 2x20: $\text{resp}(50\%) = \frac{2x}{1-(.5*.5)} = 2.66x$
 - MM1: 1x40: $\text{resp}(50\%) = x / (1-.5) = 2x$
 - MM1: 2x20: $\text{resp}(90\%) = \frac{2x}{1-.9*.9} = 10.52x$
 - MM1: 1x40: $\text{resp}(90\%) = x / (1-.9) = 10x$
 - MM1: 2x20: $\text{resp}(95\%) = \frac{2x}{1-.95*.95} = 20.51*x$
 - MM1: 1x40: $\text{resp}(95\%) = x / (1-.95) = 20x$

Processor Performance Concepts

- Which provides better response time for SINGLE THREAD?
 - 2 x 20 MIP engines, (**LESS QUEUEING**)
 - 1 40 MIP engine? (**SMALLER SERVICE TIME**)
- at what percent utilization? (“x” is cpu requirement)
 - MM1: 2x20: resp = 2x
 - MM1: 1x40: resp = x
- Common error made by sales people and financial people
- The 2x20 is less expensive in hardware and software
 - Fine for high multitasking,
 - bad for single thread or batch
- What problem is to be resolved?

Common Reported CPU Performance Problems

Problems:

- Workload timing out,
- Application running slow
- Workload/Server in CPU Wait

Causes - Tuning

- LPAR Weights vs utilization
- LPAR VCPU vs SHARE
- Share settings poor
- Operating on GP, not IFL
- Processor Utilization high

Causes - Workload

- Master processor
- Cron jobs synchronized
- Spin locks (DIAG 9C)

Processor Distribution Management

Objective: operate at high utilization

- Requires management decisions, prioritization
- Alternative to management is more hardware/licensing

Managing distribution:

- z/VM LPAR share of IFLs
 - Based on weight of LPAR
 - Based on number of IFLs (1 40 mip vs 2 20 mip)
- Linux server's share is
 - Share of z/VM LPAR (relative/absolute)
 - Based on number of vcpu
- Process “niced”, “priority”

LPAR Configuration

z/VM share of IFLs

Report: ESALPARS Logical Partition Summary

Time	CPUs	Slice	Name	<--Complex--> <-----Logical Partition----->			<-Assigned <---LPAR-->		
				Phys Dispatch	Virt CPU	<%Assigned>	Weight	Pct	
				Nbr	CPUs	Type	Total	Ovhd	
00:15:00	23	Dynamic	Totals:	0	22	CP	506.0	4.5	999 100
			Totals:	0	23	IFL	903.1	8.6	1000 100
			ZVMQA	11	6	IFL	374.8	0.9	150 15.0
			MVSPRD	7	10	CP	320.1	3.2	860 86.1
			MVSQA	1	6	CP	181.8	1.1	71 7.1
			ZVMD EQ	9	4	IFL	131.6	2.0	100 10.0
			ZVMPRD	8	10	IFL	333.7	4.9	650 65.0
			ZVMSHR	12	3	IFL	63.0	0.8	80 8.0
			MVSTST	17	3	CP	5.1	0.1	8 0.8

Totals by Processor type:

<-----CPU-----> <-Shared Processor busy->

Type	Count	Ded	shared	Total	Logical	Ovhd	Mgmt
CP	7	0	7	511.9	501.5	4.5	5.9
IFL	10	0	10	915.6	894.5	8.6	12.5
ZIIP	3	0	3	23.9	22.3	0.4	1.2

Processor Utilization Components

LPAR Physical Overhead

LPAR Assigned Logical Partition

- Logical Partition Logical Overhead
- LPAR logical time

z/VM (LPAR logical time)

- System Time (z/VM Control Program)
- Emulation (z/VM Guest time)
- User Overhead (allocated system time)

Linux (Emulation time)

- System time (kernel time)
- IRQ Time
- User time

IDLE

z/VM share of IFLs

Report: ESALPARS Logical Partition Summary

Time	CPUs	Slice	Name	<--Complex--> <-----Logical Partition----->			<-Assigned <---LPAR-->	Weight	Pct
				Phys Dispatch	Virt CPU	<%Assigned>			
00:15:00	23	Dynamic	Totals:	0	22	CP	506.0	4.5	999 100
			Totals:	0	23	IFL	903.1	8.6	1000 100
			ZVMQA	11	6	IFL	374.8	0.9	150 15.0
			ZVMDEQ	9	4	IFL	131.6	2.0	100 10.0
			ZVMPRD	8	10	IFL	333.7	4.9	650 65.0
			ZVMSHR	12	3	IFL	63.0	0.8	80 8.0

Totals by Processor type:

<-----CPU-----> <-Shared Processor busy->			Type Count Ded shared Total Logical Ovhd Mgmt			
IFL	10	0	10	915.6	894.5	8.6 12.5

- ZVMQA is allocated 150/1000 of 10 SHARED IFLs
- ZVMQA is using 37.5% of 10 SHARED IFLs
- IFLs running 91.6% busy

Each LPAR gets a weight

LPAR's share:

- $(\text{LPAR Weight}) / \text{SUM}(\text{LPAR Weights})$

Processor share of system:

- $(\text{LPAR Share}) / (\text{Number CPUs in LPAR})$

Processor share of a CPU is

- $(\text{Processor share of system}) * (\text{Number physical processors})$

LPAR Weights Example

- **ESALPAR (Partial report)**
- **Note each vcpu running at 10%?**
- **z/VM can dispatch 8 concurrent virtual machines**
 - Less queueing, slower service
 - But, each single vcpu runs “slow”

Time	<--Complex-->		<--Logical->		<-----Logical Processor----->						
	Phys CPUs	Dispatch Slice	<-Partition>	VCPU Name	<%Assigned>	Total No.	Ovhd Addr	Weight	Cap-ped	Wait Comp	
Average:	8	Dynamic	ZVM		6	0	8.3	0.2	10	No	No
				1	10.2	0.2	10	10	No	No	
				2	11.0	0.2	10	10	No	No	
				3	11.1	0.2	10	10	No	No	
				4	10.5	0.2	10	10	No	No	
				5	10.5	0.2	10	10	No	No	
				6	10.5	0.2	10	10	No	No	
				7	10.6	0.2	10	10	No	No	
				LPAR	82.8	1.4					

LPAR Share Example

Processor Details

- 30 LPARs configured
- 4 LPARs active
- Total of all active lpar shares: 60
- z/VM Weight: 10 (out of 60)
- z/VM Logical Processors: 8
- Physical processors online: 8

Guaranteed processor share (speed)

- $(\text{Share of system} / \text{nbr logical processors}) * \text{nbr phys}$
- $((10 / 60) / 8) * 8 = .16$

Each virtual cpu at peak runs at 16% rated speed

- **(go back to processor performance concepts)**

Processor Details: If change to 4 logical processors:

- 4 LPARs active
- Total of all shares: 60
- z/VM Weight: 10 (out of 60)
- z/VM Logical Processors: 4
- Physical processors online: 8

Guaranteed processor share (speed)

- $((10 / 60) / 4) * 8 = .32$
- Real problem in many installations

Too many logical processors will slow you down!

- Specifically the master processor....
- The same concept applies to Linux virtual processors
- SEE DISCUSSION ON HIPERDISPATCH!!!

LPAR Summary Report

Report: ESALPARS Logical Partition Summary TEST MAP
Monitor initialized: 08/04/03 at 18:52:10 on 2084 serial 4B54A First record

<--Complex--> <-----Logical Partition---->				<-Assigned Shares----->						
Time	Phys Dispatch		Virt <%Assigned>	<---LPAR-->		<VCPU Pct>				
	CPUs	Slice Name	Nbr	CPUs	Total	Ovhd	Weight	Pct	/SYS	/CPU
Average:	8	Dynamic Totals:	0	22	188.7	2.1	60	100		
		ZVM	6	8	82.8	1.4	10	16.0	2.00	16.0
		CF01	1	1	99.9	0.0	10	16.0	16.0	128
		LINUXSW	2	2	0	0	10	16.0	8.00	64.0
		S01	3	4	4.6	0.4	10	16.0	4.00	32.0
		S02	4	0						
		VMTPC	5	5	1.2	0.2	10	16.0	3.00	24.0
		ZVMCSS1	16	2	0.2	0.0	10	16.0	8.00	64.0

“ZVM” Allocated 16% of 8 CPUs

Each virtual cpu allocated 2% of system (8 CPUs)

Each processor rated at 16% speed of real processor

LPAR with HiperDispatch

Stated Purpose of HiperDispatch and vertical scheduler:

- Localize work to L1/L2 cache
- Reduce impact of installation configuration errors

Impact

- Virtual CPUs disabled, share redistributed
- Faster master processor
- L1/L2 cache impact negligible

ESAOPER:

```
07:00:41 CPU Park from 15 to 13 CPUUtil= "12.9",
07:00:43 CPU Unpark from 13 to 15 CPUUtil= "12.5
07:05:35 CPU Park from 15 to 13 CPUUtil= "12.2",
07:05:37 CPU Unpark from 13 to 15 CPUUtil= "12.0
07:05:53 CPU Park from 15 to 12 CPUUtil= "12.0"
07:05:55 CPU Unpark from 12 to 15 CPUUtil= "10.4
07:07:13 CPU Park from 15 to 13 CPUUtil= "12.5",
07:07:15 CPU Unpark from 13 to 15 CPUUtil= "11.9
07:07:19 CPU Park from 15 to 13 CPUUtil= "12.1",
07:07:21 CPU Unpark from 13 to 15 CPUUtil= "11.8
07:07:29 CPU Park from 15 to 13 CPUUtil= "12.1",
```

LPAR with HiperDispatch

HiperDispatch requires Vertical scheduling

```
Report: ESALPAR      Logical Partition Analysis
Monitor initialized: 05/31/16 at 00:00:00 on 2827 serial 2F5A7
-----
          <--Complex--> <--Logical-> <-----Logical Processor-
          Phys Dispatch <-Partition> VCPU <%Assigned> VCPU
Time    CPUs   Slice Name   No. Addr Total Ovhd TYPE Weight
-----  -----  -----  -----  -----  -----  -----  -----  -----
07:15:00  19  Dynamic R3SYSG   2     0  87.9   0.8 IFL   400
          1  89.3   0.7 IFL   400
.....
          11 81.9   0.8 IFL   400
          12 77.5   1.0 IFL   400
          13 75.5   0.9 IFL   400
          14 60.9   0.7 IFL   400
          LPAR 1245  11.8
          R3SYS1  3     0  48.6   2.1 IFL   500
          1  35.5   1.6 IFL   500
          2  40.4   1.7 IFL   500
          3  38.9   1.5 IFL   500
          4  36.8   1.7 IFL   500
          5  38.8   1.7 IFL   500
          6  40.1   1.3 IFL   500
          7  32.5   1.3 IFL   500
          8  30.0   1.2 IFL   500
          9  18.6   0.9 IFL   500
          10 17.8   1.4 IFL   500
          11 0.0    0.0 IFL   500
          12 0.0    0.0 IFL   500
          13 0.0    0.0 IFL   500
          14 0.0    0.0 IFL   500
          LPAR 378.1  16.4
```

Time Slice: Dynamic, used exclusively

Weights: Sets priority between Logical Partitions

Virtual processors

Capping

- Limits Assigned time to LPAR
- Useful for outsourcing, fixed contracts

Wait Completion

- “no” gives up processor if idle (default)
- “yes”, Partition keeps processor even if idle (rarely/never used)

CEC Processor Utilization:

- Physical overhead
- Assigned time (physical cpu assigned to logical cpu)
 - Logical overhead
 - Assigned time (work)

Each LPAR gets a “Share”

LPAR was based on CP, much of LPAR has the same concepts as CP

Physical overhead is cost of dispatching Ipar vcpus

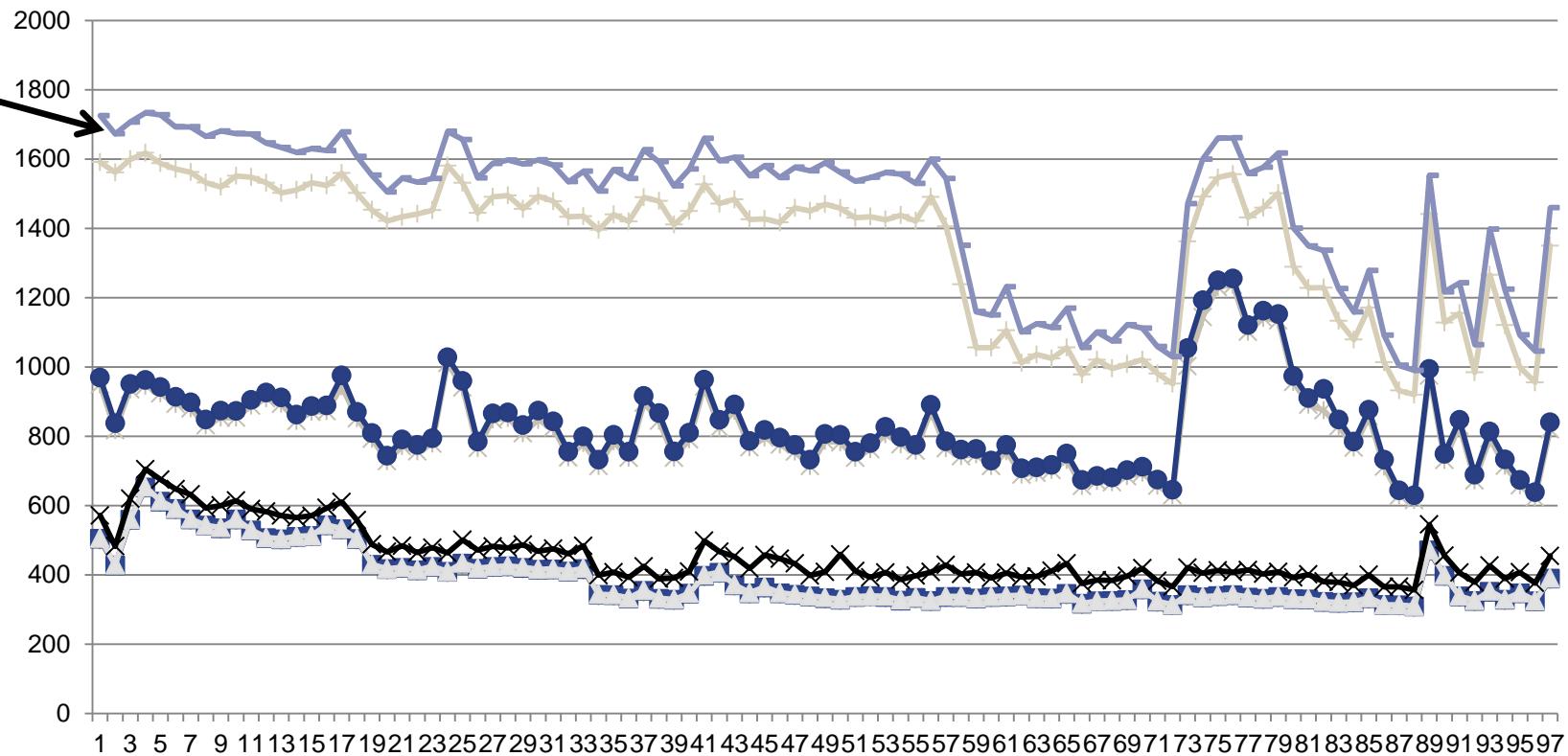
Linux Servers

- 120 servers total (Big, ORACLE)
 - 4gb-40gb
 - (1 / 2 size from original SUN servers)

Hardware

- 17 IFLs
 - 7 servers per IFL
 - 395 vcpus (23:1 overcommit)
- 7 LPARs, each with 17 VCPU
 - Worst case possible for overhead

LPAR Configuration Overhead

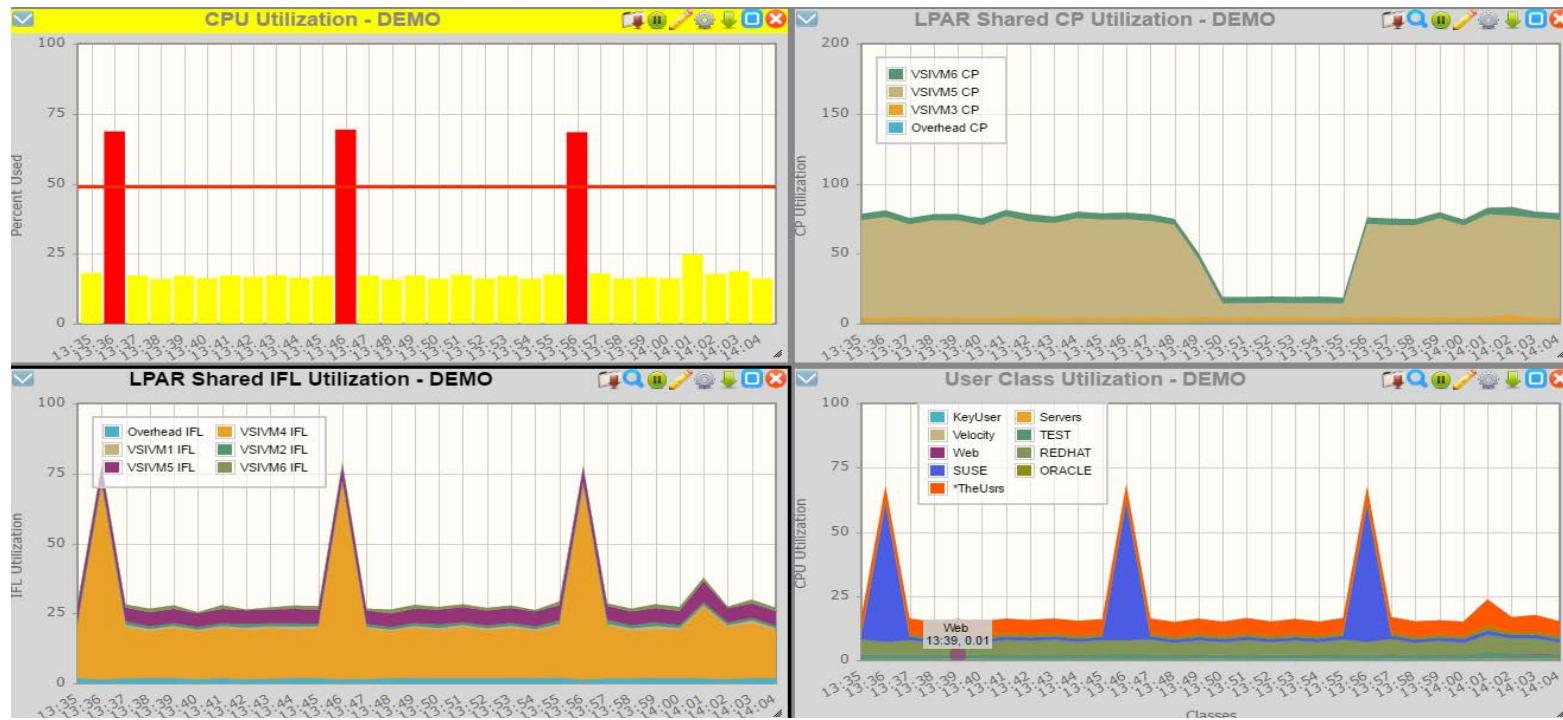


17 IFLs, 7 lpars, 17 vcpus each, **7:1 overcommit**

Physical Overhead significant from real processor overcommit

Analyzing IFL Consumption

- zVIEW, 4 charts (default subset)
- IFL utilization, CP utilization
- z/VM LPAR utilization, by workload (classes)



Managing Distribution - What is my share?

CPU distribution depends on

- **LPAR Weight (and impact of other LPARs)**
- **Share within z/VM LPAR**
- **Impact of other virtual servers**

Share analysis:

- **ESALPARS for LPAR / IFL utilization**
- **ESALPARS for LPAR allocation**
- **ESACPUU / ESACPUA for z/VM perspective**
- **ESAUSRC for share settings**
- **ESAUSP2 for how much am I getting out of used**

“z” Processor Overview (ESAHDR)

Machine Model/Type	z9EC:2094/707
System Sequence Code	000000000006EA7D
Processor 0 model/serial	2094-707 /0AEA7D Master
Processor 1 model/serial	2094-707 /0AEA7D
Processor 2 model/serial	2094-707 /0AEA7D
Processor 3 model/serial	2094-707 /0AEA7D

ESAME (Memory Extension) Nucleus in use

Power of processor in terms of service Units: **25000**

ESA/370 hardware installed

Operating on IFL Processor(s)

Channel Path Measurement Facility(CPMF) Extended is installed

Service units from table
Understand the
CEC (two books)
z/VM (IFLs)

Processor Measurement

Report: ESACPUU				CPU Utilization Report							Linux Test			
Monitor initialized: 05/06/08 at 12:00:00 on 2094 serial AEA7D											First record analyzed:			
Time	<----Load---->			<----CPU (percentages)---->				<----External (per second)----						
	<-Users-> Tran			Total	Emul	User	Sys	Idle	<--Page-->		<--Spool-->		RSCH+	
Time	Actv	In Q	/sec	CPU	util	time	ovrhd	ovrhd	time	Read	Write	Read	Write	SSCH
12:01:00	103	118	9.1	0	92.8	88.6	2.3	1.9	7.2	11	52	0	0	220
				1	93.8	90.5	2.2	1.0	6.2	14	0	0	0	182
				2	94.4	90.9	2.2	1.2	5.6	17	0	0	0	196
				3	94.5	90.9	2.1	1.5	5.5	13	0	0	0	179
System:				375.4	361.0	8.9	5.5	24.4	55	52	0	0	0	778

- Processor utilization has three components:
 - Emulation time – running users in Interpretive Execution
 - User overhead – CP time performing services for a user
 - System overhead – CP “housekeeping”
 - Note master processor – only problem if architecturally constrained

Processor Measurements User View

ESAUSP2: CPU Consumption in percent

- Total all user
- By user
- By Class

Note

- one server
dominates CPU
- Capture ratio

T:V Ratio is Total to
Virtual, 1.0 is good

```
Report: ESAUSP2      User Resource Rate Report
Monitor initialized: 05/06/08 at 12:00:00 on 2094 serial
-----
          <---CPU time--> <----Main Storage (pages)---->
UserID  <(Percent)> T:V <Resident> Lock <----WSS----->
/Class   Total   Virt  Rat Totl Activ -ed Totl Activ Avg
-----
12:01:00 369.9 361.0 1.0   17M   17M  417  17M   17M 129K
***User Class Analysis***
*Servers 1.95  1.72 1.1 7566  7555  49  8674  7444  207
*CMSUser 0.00  0.00 23   0     0     0    63    63   63
*Linux   184.0 180.6 1.0 15M   15M  305  15M   15M 185K
*Misc    183.7 178.5 1.0 2M   1642K 11   2M   1642K 328K
***Top User Analysis***
LXPWK001 183.5 178.4 1.0 2M   1641K 3   2M   1641K 2M
LXWKB215 37.63 37.01 1.0 782K  782K  1   782K  782K 782K
LXWKB211 33.97 33.88 1.0 514K  514K  0   514K  514K 514K
LXWKB210 17.64 17.55 1.0 298K  298K  2   298K  298K 298K
LXWKB214 16.86 16.68 1.0 1M   1188K 0   1M   1254K 1M
LXWKB228 6.01  5.98 1.0 731K  731K  3   731K  731K 731K
LXWKB222 5.06  4.94 1.0 621K  621K  5   621K  621K 621K
LXWKB183 4.70  4.57 1.0 231K  231K  0   230K  230K 230K
LXWKB220 3.69  3.66 1.0 125K  125K  8   124K  124K 124K
LXWKB225 3.65  3.52 1.0 780K  780K  0   780K  780K 780K
ESATCP   0.45  0.35 1.3 1038  1038  1   1037  1037 1037
TCPIP2   0.02  0.01 2.0 1142  1142  48  198   198 198
```

Much system code NOT re-entrant

- Must be single threaded
- Can not update one control block by multiple processors simultaneously

Implementation

- hardware locks: TS, CS, CDS instructions
- software locks: “ownership” of resources
 - (such as in database)
- running on the Master Processor

SPIN Locks

- Test for lock, if fail, test for lock
- Linux uses “spin lock”, replaced with Diag44 -> DIAG9C
- Linux spin locks an issue, cost in CPU

Resource Serialization Master Processor

Many CP processes run “master only” to ensure the integrity of the system

- Spooling
- some IUCV services (*MSG, *RPI, *ACCOUNT from CP)
- Page migration
- execution of ALL CP commands
- Line mode console I/O

Master processor utilization shows up as higher System Overhead and User Overhead on the Master Processor.

Higher Master CPU busy is worse on a system with more processors.

- Master calls is measured
- Simulation wait is measured
- Processor imbalance can be a problem

Master Processor Problem

CPU Example

- User overhead high on master
- System overhead high on master
- Master processor can be a limiter

Report: ESACPUU CPU Utilization													
Time	<----Load---->		<----CPU (percentages)---->				<----External (per second)---->						
	<-Users->	Tran	Total CPU	util	Emul ovrhd	User ovrhd	Sys	Idle	<--Page-->	<--Spool-->	RSCH+	SSCH	ExInt
09:19:12	7	5.0	0.1	1	99.4	20.9	58.8	19.8	0	0	0	0	3 140
				2	84.7	43.6	30.7	10.3	15.0	0	0	0	0 154
				3	84.2	43.2	30.9	10.1	15.5	0	0	0	0 153
				4	84.5	43.6	31.1	9.7	15.2	0	0	0	0 155
System:					352.7	151.3	151.6	49.9	45.7	0	0	0	3 602

Would adding another processor help this system?

MASTER Processor Analysis

- **Investigating Master Overhead**
 - User overhead high on master
 - System overhead high on master
 - All spool I/O performed by master Note I/O lower on master?
- **ESAIUCV: RPI, MSG IUCV**
- **ESACPUA: Page migration, Spool I/O**

Time	<----Load--->				<-----CPU (percentages)----->				<--External (per second)-->						
	<-Users>	Tran	Total	Emul	User	Sys	Idle	.	<--Page-->	<-Spool-->	RSCH+	SSCH			
	Actv	InQ	/sec	CPU	util	time	ovrhd	ovrhd	time	.	Read	Write	Read	Write	
10:21:09	255	20	23.9	0	86.4	16.7	53.6	16.2	13.4	.	60	71	9	4	86
				1	67.9	35.9	23.7	8.4	31.9	.	82	56	0	0	135
				2	69.5	36.4	24.5	8.7	30.3	.	89	58	0	0	140
				3	67.0	35.4	23.2	8.4	32.8	.	76	71	0	0	134
										.					
System					290.9	124.3	124.9	41.7	108.4	.	307	256	9	4	496

Processor Local Dispatch Vectors

- **Users are selected for dispatch from a PLDV.**
- **Each processor has a local PLDV.**
- **Master ONLY PLDV**
 - Some work for users
 - Some System work
 - The System VMDBK
- **The System VMDBK and users on the Master Only PLDV are dispatched on the Master Processor.**

Dispatch Vector Activity

Report: ESAPLDV				Processor Local Dispatch Vector Activity				Linux Test		ESAMAP 3.7.4			
Time	<----Users---->			Tran		<VMDBK Moves/sec>		<-----PLDV Lengths----->				Dispatcher	
	Logged	Actv	In Q	/sec	CPU	Steals	To Master	Avg	Max	Mstr	MstrMa	%Empty	Long Paths
12:01:00	129	103	118	9.1	0	0	2.5	3.2	4.0	0.0	1.	8.3	4497.1
					-	1	0	0	2.1	4.0	.	38.3	3942.1
						2	0	0	2.0	4.0	.	41.7	3942.7
						3	0	0	1.8	3.0	.	38.3	3741.7
System:							0	2.5	9.2	15.0	0.0	1.	126.7
													16123.5

- The Dispatcher selects users from the CPU's Processor Local Dispatch Vector (PLDV).
- Each processor has a Local PLDV.
- The Master Processor has a special PLDV from which “master only” work for users is selected.

Processor Utilization (Is it correct?)

- **Why is the Linux LPAR always 100% busy? (RMF)**
 - Dedicated processors show up as 100%
- **Processor utilization measurements are NOT consistent between products/commands**
 - This matters when running second level, or in an LPAR
 - CP Indicate shows percent of what is available
 - CP Does NOT KNOW what is available
- **ESAMON/ESAMAP show absolute utilization**
 - LPAR information is available
- **Measuring Linux under VM or LPAR had same issue**
Linux reported a percent of a percent busy

Effects of Logical Partitioning

Report: **ESASSUM Subsystem** Activity Velocity Software, Inc.

Time	<---Users----> Transactions <Processor> Storage (MB) <-Paging--> <----I/O-----> <MiniDisk> Spool												<-avg number-> Per Avg. Utilization Fixed Active <pages/sec> <-DASD--> Other <-Cache--> Page				
	On	Actv	In Q	Minute					User	Resid.	XStore	DASD	Rate	Rate	Rate	%Hit	Rate
08:00:08	1479	244	34.3	1310.1	0.603	124	87	36.9	192.0	888	451	641	15.4	40	687.9	49.3	36
08:01:08	1500	248	46.0	1260.9	0.543	147	110	37.3	192.7	904	494	732	20.1	37	881.6	53.9	32
*****Summary*****																	
Average:	1483	245	37.3	1297.8	0.589	130	93	37.0	192.1	892	461	664	16.7	39	736.4	50.7	35

- A high-level view of processor utilization shows a system with some capacity to spare.
- Next step, look at processor configuration

Effects of Logical Partitioning

Report: ESACPUU CPU Interval Analysis Velocity Software, Inc.

Time	<----Load---->			<----CPU (percentages)---->					<---Internal (per second)---->				
	<-Users-> Tran			Total	Emul	User	Sys	Idle	Diag-	Inst.	SIE	Fast	Page
	Actv	In Q	/sec	CPU	util	time	ovrhd	ovrhd	time	nose	sim.	intrcp	path fault
08:00:08	244	34.3	24.6	0	48.5	27.3	16.7	4.5	9.9	1449	1478	1753	0 18
				1	35.9	28.8	5.2	1.9	11.8	818	599	716	0 9
				2	39.5	31.4	5.9	2.2	13.5	902	682	815	0 11
System:				124.0	87.4	27.9	8.7	35.3	3170	2758	3284	0	37
08:01:08	248	46.0	24.0	0	53.6	32.5	16.7	4.4	7.1	1557	1588	1806	0 24
				1	44.6	37.2	5.4	1.9	6.5	843	594	685	0 11
				2	48.8	40.2	6.4	2.2	7.4	903	704	817	0 12
System:				147.0	109.9	28.5	8.6	21.0	3303	2886	3308	0	48

- A more detailed view of processor utilization seems to confirm this hypothesis. CPU to spare.

Effects of Logical Partitioning

Report: ESAXACT Transaction Analysis Velocity Software, Inc.

		Percent non-dormant																
User ID	<-Samples->	Total	In Q	Run	Sim	CPU	SIO	Pg	E-	D-	T-	Tst	Lim	Pct				
/Class									SVM	SVM	SVM	CF	Idl	I/O	Ldg	Oth	Lst	Elig
System:	5936	149	5.4	34	8.7	0	3	0	0	6.0	2	36	4.7	.	0	.	0	
Hi-Freq:	176K	7057	2.0	17	2.8	0	1	0	3.8	4.2	49	17	3.1	0	0	.	0	
***Resource use by User Class																		
*Servers	3720	568	3.0	29	4.2	0	0	0	21	6.9	1	28	7.6	0	0	.	0	
*Keys	1080	490	1.6	0.6	6.7	0	0	0	16	19	1	43	13	0	0	.	0	
*TheUsrs	172K	6108	1.9	16	2.6	0	1	0	1.2	3.0	57	14	2.5	0	0	.	0	

- User state sampling shows significant amount of CPU wait as compared to “Running”, Simulation wait even greater.

Effects of Logical Partitioning

Report: ESALPAR Logical Partition Analysis Velocity Software, Inc.

<----Load----> <--Complex--> <--Logical-> <-----Logical Processor----->													
<-Users->		Tran	Phys	Dispatch	<-Partition>	VCPUs	<%Assigned>	Cap-	Wait				
		Slice	Name	No.	Addr	Total	Ovhd	Weight	ped Comp				
08:00:08	244	34.3	24.6	3	Dynamic CMS2	1	0	58.7	0.2	155	No	Yes	
						1	47.8	0.1		155	No	Yes	
						2	53.2	0.1		155	No	Yes	

						LPAR 159.7 0.4							
					SWCF	2	0	36.6	0.1	130	No	Yes	
						1	43.0	0.1		130	No	Yes	
						2	46.7	0.1		130	No	Yes	

						LPAR 126.3 0.3							
					CMS8	3	0	9.1	0.1	15	No	Yes	
						1	4.6	0.2		15	No	Yes	

						LPAR 13.7 0.3							
Total Logical Partition busy:										299.6			
Total Physical Management time: 0.366													

- This system does not have access to 100% of each processor.
- Reducing CMS2 LPAR to 2 processors will perform better.

Processor Capture Ratio

Report: ESASSUM

Subsystem Activity

Time	<---Users---->			Transactions		<Processor>		Capture Ratio
	<-avg number->			Per	Avg.	Utilization		
	On	Actv	In Q	Minute	Resp	Total	Virt.	
16:26:00	741	154	27.0	765.0	0.521	799	678	99.96
16:27:00	738	136	41.0	602.0	0.356	693	585	99.99
16:28:00	740	131	31.0	695.0	0.318	688	589	99.99
16:29:00	739	159	32.0	665.0	0.437	858	735	100.09
16:30:00	738	154	34.0	783.0	0.435	656	558	100.04

Average:	738	157	35.9	793.6	0.416	740	632	99.50

Capture ratio is CPU accounted for / CPU used

Capture ratio critical for capacity planning, chargeback

Unique methodology

Managing the delivery of CPU Resource

Options for tuning processor

- **Reducing overhead**
 - Application tuning to reduce processor demand
 - Reducing system overhead
- **Reallocating resources**
 - Adjust SHARE values
 - Limit DSPBUF occupancy
 - Dedicate processors
 - Interactive BIAS (SET SRM IABIAS)

Limiting Shares

```
Q share vmservu
```

```
USER VMSERVU :RELATIVE SHARE= 100 MAXIMUM SHARE= NOLIMIT
Ready; T=0.01/0.01 16:58:54
```

LIMITs

- LIMITHARD caps resource consumption regardless of other user demands
- LIMITSOFT caps resource consumption unless all users have received their target minimum, and there are no unlimited users who can consume resources

```
set share vmservu relative 200 500 limitsoft
```

```
USER VMSERVU : RELATIVE SHARE= 200 MAXIMUM SHARE=LIMITSOFT
RELATIVE 500
```

```
Ready; T=0.01/0.01 17:01:12
```

```
set share mvsy়1 abs 5% abs 20% limithard
```

```
USER MVSYS1 : ABSOLUTE SHARE = 5%
```

```
MAXIMUM SHARE = LIMITHARD ABSOLUTE 20%
```

```
Ready; T=0.01/0.01 14:40:49
```

Adjusting DSPBUF values

- The purpose of the Dispatch Buffer (DSPBUF) is the limit on the number of users of each Scheduler class allowed in the Dispatch List – thereby controlling CPU utilization
- If CPU is 100% busy, try lowering the Q3 value.
 - Watch CPU busy very carefully.
- What happens if DSPBUF is lowered too far for Q3?
 - (Hint: How is Elapsed Time Slice determined?)
- Experimenting with DSPBUF is not for the casual tuner or the faint of heart.

I/O Processor is another type of processor

If I/O processor overloaded, determine why.

- Is I/O Slow with no explanation?
- Could be iop....
- ESCON and Ficon mixed channels?
- DASD Cache performing poorly?

Report: Report: ESAIOP		I/O Processor Analysis							
Time	Nmbr	I/O <-----I/O Processor----->				<-Percent of Strts busy->			
		Proc	<Pct Util>	<Rate/Second>	chan	switch	CtlUnit	Device	
21:22:00	0	65.8	34.2	2690	2402	1488	12.7	0.3	1.5
	1	55.6	44.4	2939	3686	837	4.3	0.3	0.1
	2	54.2	45.8	1684	2979	1756	0.6	0.5	0.4
21:27:00	0	50.6	49.4	2227	2190	1374	12.6	0.4	0.9
	1	39.8	60.2	2540	2997	648	3.6	0.4	0.2
	2	41.2	58.8	1711	2691	1334	0.7	0.4	0.5

I/O Processor is another type of processor

If I/O processor overloaded, determine why.

- Is I/O Slow with no explanation?
- Could be iop....
- ESCON and Ficon mixed channels?
- DASD Cache performing poorly?

Report: Report: ESAIOP I/O Processor Analysis

Report: ESAUSRQ User Queue and Load Analysis Veloc
Monitor initialized: 05/22/08 at 14:00:00 on 2084 serial 15BAF First

<-----User Load----->				<-----Average Num----->						
UserID /Class	Logged on	Non- Idle	Active	Disc- conn	Total InQue	Tran /min	Q0	Q1	Q2	Q3
14:01:00	1061.0	.	156.0	.	20.0	1175	5.0	6.0	2.0	7.0
14:02:00	1063.0	.	157.0	.	25.0	1184	10.0	5.0	4.0	6.0
14:03:00	1064.0	.	188.0	.	52.0	1423	3.0	2.0	8.0	39.0
14:04:00	1062.0	.	203.0	.	38.0	1245	3.0	1.0	6.0	28.0
14:05:00	1061.0	.	193.0	.	46.0	1228	5.0	8.0	4.0	29.0
.....										

Processor Case Study

User's complain, InQueue skyrockets, why?

- Impact really is quickdsp and Q3 –Really long running transactions.

Report: ESAUSRQ		User Queue and Load Analysis							Veloc	
UserID /Class	Logged on	<-----User Load----->			<-----Average Num Dispatch List--					
		Non- Idle	Active	Disc- conn	Total InQue	Tran /min	Q0	Q1	Q3	
<hr/>										
14:01:00	1061.0	.	156.0	.	20.0	1175	5.0	6.0	2.0	7.0
14:02:00	1063.0	.	157.0	.	25.0	1184	10.0	5.0	4.0	6.0
14:03:00	1064.0	.	188.0	.	52.0	1423	3.0	2.0	8.0	39.0
<hr/>										
14:18:00	1064.0	.	154.0	.	31.0	1185	10.0	5.0	6.0	10.0
14:19:00	1065.0	.	161.0	.	36.0	1130	6.0	4.0	3.0	23.0
14:20:00	1065.0	.	186.0	.	47.0	1143	13.0	3.0	1.0	30.0
14:21:00	1066.0	.	190.0	.	72.0	1140	25.0	15.0	8.0	24.0
14:22:00	1065.0	.	213.0	.	73.0	1189	35.0	3.0	1.0	34.0
14:23:00	1067.0	.	243.0	.	88.0	1157	31.0	3.0	1.0	53.0
14:24:00	1067.0	.	259.0	.	81.0	1105	11.0	2.0	5.0	63.0
14:25:00	1067.0	.	215.0	.	46.0	932	12.0	6.0	3.0	25.0
<hr/>										
14:30:00	1069.0	.	266.0	.	108.0	1227	34.0	6.0	7.0	61.0
14:31:00	1069.0	.	274.0	.	116.0	1183	30.0	2.0	2.0	82.0
14:32:00	1067.0	.	266.0	.	126.0	960	47.0	4.0	5.0	70.0
14:33:00	1067.0	.	257.0	.	105.0	1230	43.0	7.0	13.0	42.0

Processor Case Study

Check processor, cpu is a constant, I/O is constant

Report: ESASSUM		Subsystem Activity								Velocity Software				
Time		<---Users----> Transactions <Processor> Storage (MB) <-Paging--> <----I/O---->				<-avg number-> Per Avg. Utilization Fixed Active <pages/sec> <-DASD--> Other				User Resid.	XStore	DASD Rate	Resp	Rate
		On	Actv	In Q	Minute	Resp	Total	Virt.	User Resid.					
14:01:00	1061	156	20.0	763.0	0.733	41	35	18.5	999.5	5	5	536	1.0	27.5
14:02:00	1063	157	25.0	803.0	0.594	41	35	18.5	1022.0	7	4	634	1.0	27.8
14:03:00	1064	188	52.0	981.0	1.112	41	35	18.5	1162.0	7	5	318	1.0	33.4
14:18:00	1064	154	31.0	729.0	1.055	41	36	18.5	986.5	0	3	277	1.0	26.3
14:19:00	1065	161	36.0	727.0	0.704	41	34	18.5	1061.1	226	3	303	1.3	35.3
14:20:00	1065	186	47.0	773.0	1.954	41	35	18.5	1315.9	432	2	377	1.1	30.8
14:21:00	1066	190	72.0	843.0	2.160	41	34	18.7	1308.9	1	2	769	0.8	38.9
14:22:00	1065	213	73.0	833.0	2.367	41	35	18.7	1394.9	1	3	548	0.9	31.1
14:23:00	1067	243	88.0	830.0	2.824	41	35	18.9	1537.0	1	3	858	0.8	29.8
14:24:00	1067	259	81.0	775.0	2.389	41	34	18.7	1660.4	13	3	683	0.8	18.2
14:25:00	1067	215	46.0	509.0	1.095	41	34	18.7	1452.4	8	2	583	0.8	28.5
14:30:00	1069	266	108	838.0	1.623	41	35	19.2	1618.2	5	3	511	0.8	28.8
14:31:00	1069	274	116	787.0	0.655	41	35	19.2	1630.7	8	3	569	0.8	29.0
14:32:00	1067	266	126	650.0	1.191	41	34	19.2	1580.9	4	3	774	0.8	30.7

Processor Case Study

Check LPAR Configuration

- Check weights
- VM shares with MVS and TEST, share is $179 / (179+260+5) = 40\%$
- (Only one CP defined)
- VM LPAR is capped!!!! At 40% of one CPU. VM running 100%

Report: ESALPARS		Logical Partition Summary				Velocity Software					
Time	Complex	<--Logical Partition-->			<-Assigned Shares---->			Proce	Wait	Type	
		Phys Dispatch	Virt	<%Assigned>	<---LPAR-->	<VCPU Pct>	Cap-				
Time	Complex	CPUs	Slice Name	Nbr CPUs	Total	Ovhd	Weight	/SYS	/CPU	Caped	Comp
14:01:00		1	Dynamic Totals:	0	3	80.4	0.5	444	100		
			VM	1	1	41.2	0.1	179	40.0	40.0	40.0
			MVS	2	1	39.2	0.4	260	59.1	59.1	59.1
			TEST	3	1	0	0	5	1.0	0.96	0.96
			TESTTEST	5	0						

Processor Case Study

Check User Wait States

- Running went down as percent of non-dormant, inqueue time.
- CPU wait stayed the same
- Asynchronous I/O wait is bottleneck – but DASD I/O was constant?
- Clue – something was on the Limit List – this is result of SHARE CAP
- Wait state sampling tests I/O Wait before testing Limit. If I/O wait, stops.

Report: ESAXACT		Transaction Delay Analysis													Velocity Software			
UserID	/Class	<-----Percent non-dormant----->													Times			
		<-Samples->			E-	D-	T-	Tst	<Asynch>			Lim	Pct	I/O				
Total	In Q	Run Sim	CPU	SIO	Pag	SVM	SVM	SVM	CF	Idl	I/O	Pag	Ldg	Oth	Lst	Elig	Thrott	
1061	20	5.0	5.0	40	0	0	0	0	10	0	35	0	.	0	0	0	.	
Hi-Freq:	62599	1880	3.1	1.5	39	2.8	0	0	23	4.3	3.3	22	0.8	0	0	3.0	0	0
1069	116	0.9	0.9	34	0	0	0	0	1.7	0	3.4	59	.	0	0	0	.	
Hi-Freq:	64140	7755	0.7	1.2	39	1.0	0	0	9.1	2.1	0.3	4.0	42	0	0	0.5	0	0
1067	125	0	4.0	46	0	0	0	0	2.4	0	5.6	42	.	0	0	0	.	
Hi-Freq:	64020	7508	0.8	1.2	42	1.0	0	0	8.7	2.1	0.3	3.7	40	0	0	0.5	0	0

Processor Case Study

Check User Share settings

- Cap on the database servers
- CPU consumption reaches point where database servers are limited
- Fall over the cliff
- Solution: Remove all caps. z/VM does a better job

Report: ESAUSRC User Configuration

<----- SHARE ----->							
UserID	ClassID	Account	ACI	Grp	<Normal>		<-Maximum>
					Rel	Abs	Type Share Limit
TIFSHRE	*BMAdmn	SYSTEMS	.	200	.	Abs	10.0 Soft
TIFSHRE2	*BMAdmn	SYSTEM	.	200	.	Abs	10.0 Soft
TIFSHRE3	*BMAdmn	SYSTEMS	.	200	.	Abs	10.0 Soft
TIFSHRE4	*BMAdmn	SYSTEM	.	200	.	Abs	10.0 Soft